

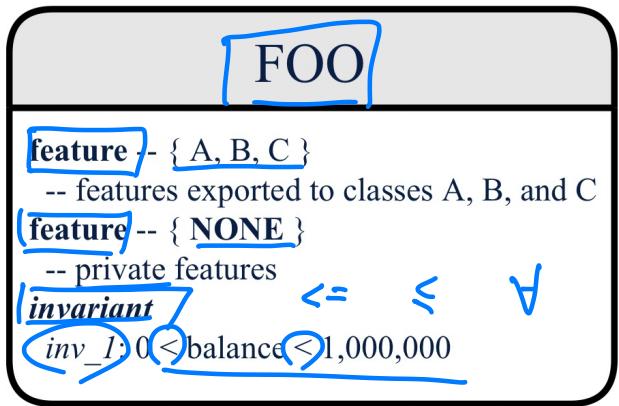
Lecture 4

Part 1

Architectural Design Diagrams

Classes: Detailed vs. Compact

Detailed View



Compact View



Contracts: Mathematical vs Programming

ARRAYED_CONTAINER + *effective*

feature -- Queries *public*

count + **INTEGER**
-- Number of items stored in the container

feature -- Commands *public*

assign *at* + **i**: **INTEGER**; **s**: **STRING**
-- Change the value at position '*i*' to '*s*'.

require

valid_index: $1 \leq i \leq \text{count}$

ensure

size_unchanged: *imp.count* = (**old imp.twin**).count

item_assigned: *imp[i]* ~ *s*

others_unchanged: $\forall j : 1 \leq j \leq \text{imp.count} : j \neq i \Rightarrow \text{imp}[j] \sim (\text{old imp.twin})[j]$

feature -- { NONE } *private*

imp + **ARRAY[STRING]**
-- Implementation of an arrayed-container

invariant

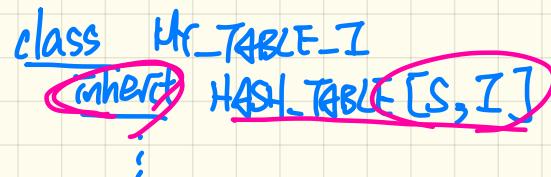
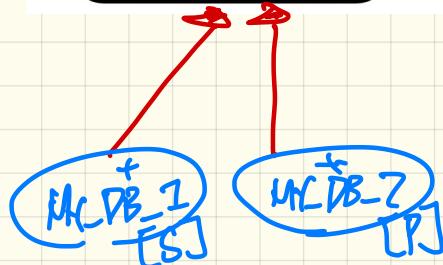
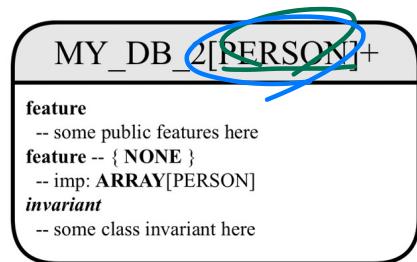
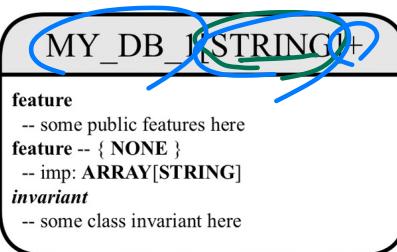
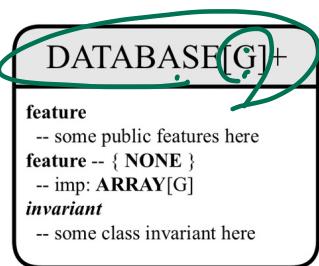
consistency: *imp.count* = *count*.

Generic Classes

- Type parameter(s) of a class may or may not be instantiated:



- If necessary, present a generic class in the detailed form:



Programming Classes: Deferred vs. Effective

```
deferred class  
  DATABASE [G]  
  feature -- Queries  
    search (g: G): BOOLEAN  
      -- Does item 'g' exist in database?  
    deferred end  
end
```

* DATABASE [G]
search (g: G): Boolean
defered

```
class  
  DATABASE_V1 [G]  
  inherit  
    DATABASE [G]  
  feature -- Queries  
    search (g: G): BOOLEAN  
      -- Perform a linear search on the database.  
    do end  
  end
```

effective

DB_V1 [G]
search +
DB_V2 [G]
search ++

```
class  
  DATABASE_V2 [G]  
  inherit  
    DATABASE_V1 [G]  
    redefine search end  
  feature -- Queries  
    search (g: G): BOOLEAN  
      -- Perform a binary search on the database.  
    do end  
  end
```

redefined

Presenting Deferred/Effective Classes in Compact Form

LIST[G]*

LINKED_LIST[G]+

ARRAYED_LIST[G]+

LIST[LIST[PERSON]]*

LINKED_LIST[INTEGER]+

ARRAYED_LIST[G]+

DATABASE[G]*

DATABASE_V1[G]+

DATABASE_V2[G]+

Presenting Deferred/Effective Classes in Detailed Form

DATABASE[G]*

feature {NONE} -- Implementation
data: ARRAY[G]

feature -- Commands
add_item*(g: G)
-- Add new item 'g' into database.
require
non_existing_item $\neg \exists$ exists(g)
ensure
size_incremented: count = old count + 1
item_added: exists(g)

feature -- Queries
count+: INTEGER
-- Number of items stored in database
ensure
correct_result: Result = data.count

exist*(g: G): BOOLEAN
-- Does item 'g' exist in database?
ensure
correct_result: Result = $(\exists i : 1 \leq i \leq \text{count} : \text{data}[i] \sim g)$

DATABASE_V1[G]+

feature {NONE} -- Implementation
data: ARRAY[G]

feature -- Commands
add_item*(g: G)
-- Append new item 'g' into end of 'data'.
feature -- Queries
count+: INTEGER
-- Number of items stored in database
exist*(g: G): BOOLEAN
-- Perform a linear search on 'data' array.

DATABASE_V2[G]+

feature {NONE} -- Implementation
data: ARRAY[G]

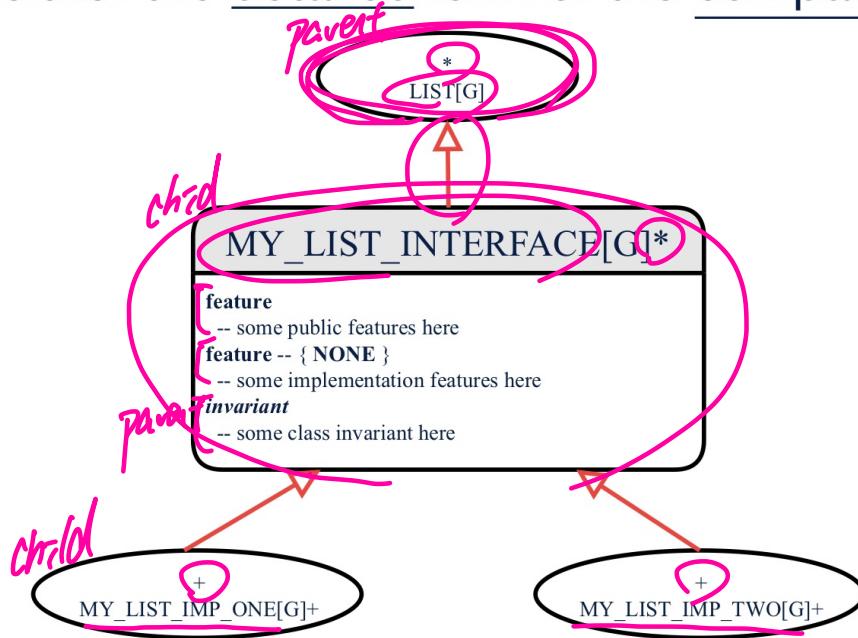
feature -- Commands
add_item++(g: G)
-- Insert new item 'g' into the right slot of 'data'.
feature -- Queries
count+: INTEGER
-- Number of items stored in database
exist\$(g: G): BOOLEAN
-- Perform a binary search on 'data' array.

invariant
sorted_data: $\forall i : 1 \leq i < \text{count} : \text{data}[i] < \text{data}[i + 1]$



Inheritace (1)

- You may choose to present each class in an inheritance hierarchy in either the detailed form or the compact form:

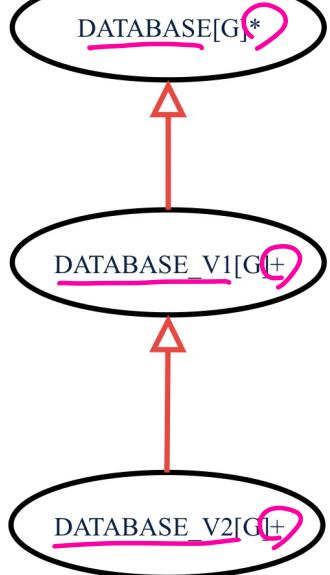


Inheritance (2)

More examples (emphasizing different aspects of DATABASE):

Inheritance Hierarchy

Features being (Re-)Implemented



detailed

DATABASE[G]*

```
feature {NONE} -- Implementation
data: ARRAY[G]

feature - Commands
  add_item*(g: G)
    -- Add new item 'g' into database.
    require
      !exists_existing_item: ~exists(g)
    ensure
      size_increased: count = old count + 1
      item_added: exists(g)

feature - Queries
  count+: INTEGER
    -- Number of items stored in database
  ensure
    correct_result: Result = data.count

exists*(g: G): BOOLEAN
  -- Does item 'g' exist in database?
  ensure
    correct_result: Result = (exists i: 1 ≤ i ≤ count : data[i] = g)
```

compact



detailed

DATABASE_V2[G]+

```
feature {NONE} -- Implementation
data: ARRAY[G]

feature - Commands
  add_item*(g: G)
    -- Insert new item 'g' into the right slot of 'data'.
    invariant
      sorted_data: ∀i: 1 ≤ i < count : data[i] < data[i + 1]

feature - Queries
  count+: INTEGER
    -- Number of items stored in database

exists+(g: G): BOOLEAN
  -- Perform a binary search on 'data' array.
```

Programming Client-Supplier Relation

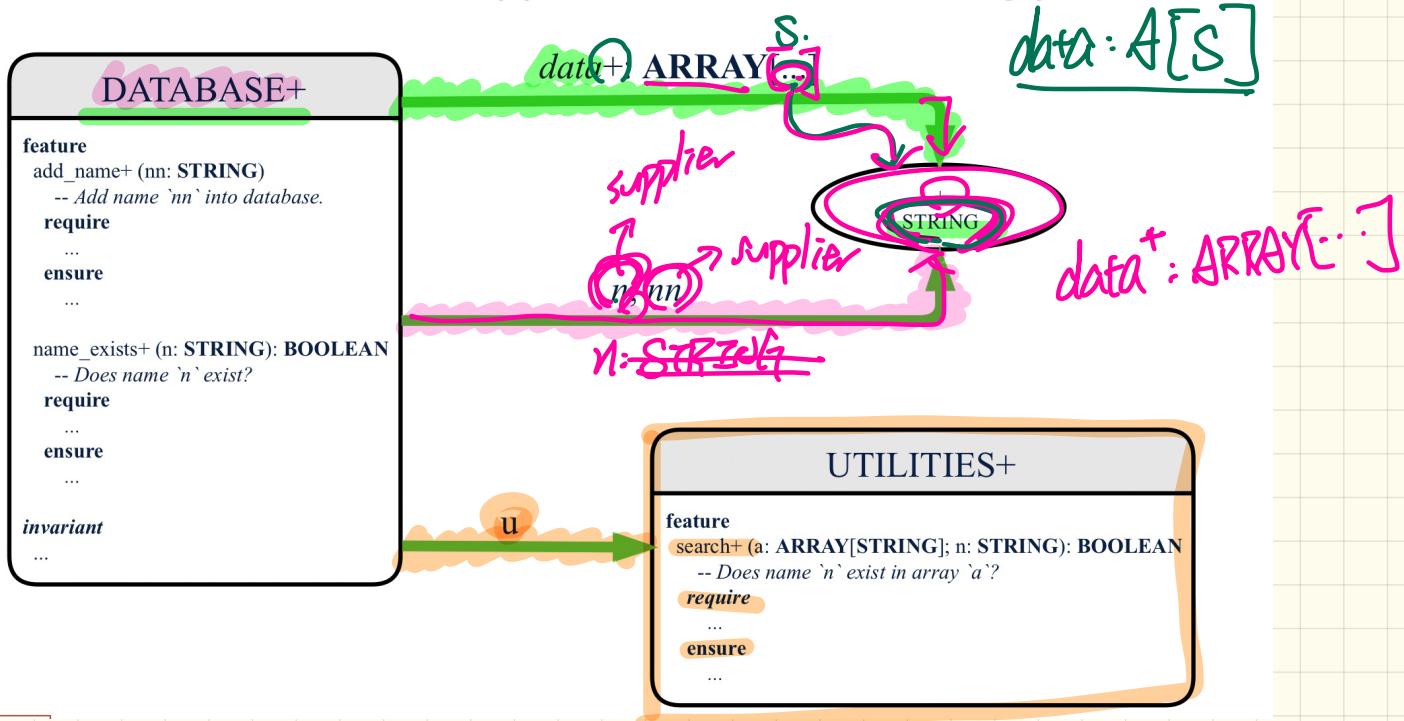
```
class DATABASE
feature {NONE} -- implementation
  data: ARRAY[STRING] → query [do]
  feature -- Commands
    add_name (nn: STRING)
      -- Add name 'nn' to database.
      require ... do ... ensure ... end

    name_exists (n: STRING): BOOLEAN
      -- Does name 'n' exist in database?
      require ...
      local
        u: UTILITIES
      do ... ensure ... end
  invariant
    ...
end
```

```
class UTILITIES
feature -- Queries
  search (a: ARRAY[STRING]; n: STRING): BOOLEAN
    -- Does name 'n' exist in array 'a'?
    require ... do ... ensure ... end
end
```

Presenting CS Relation in Diagram: Approach 1

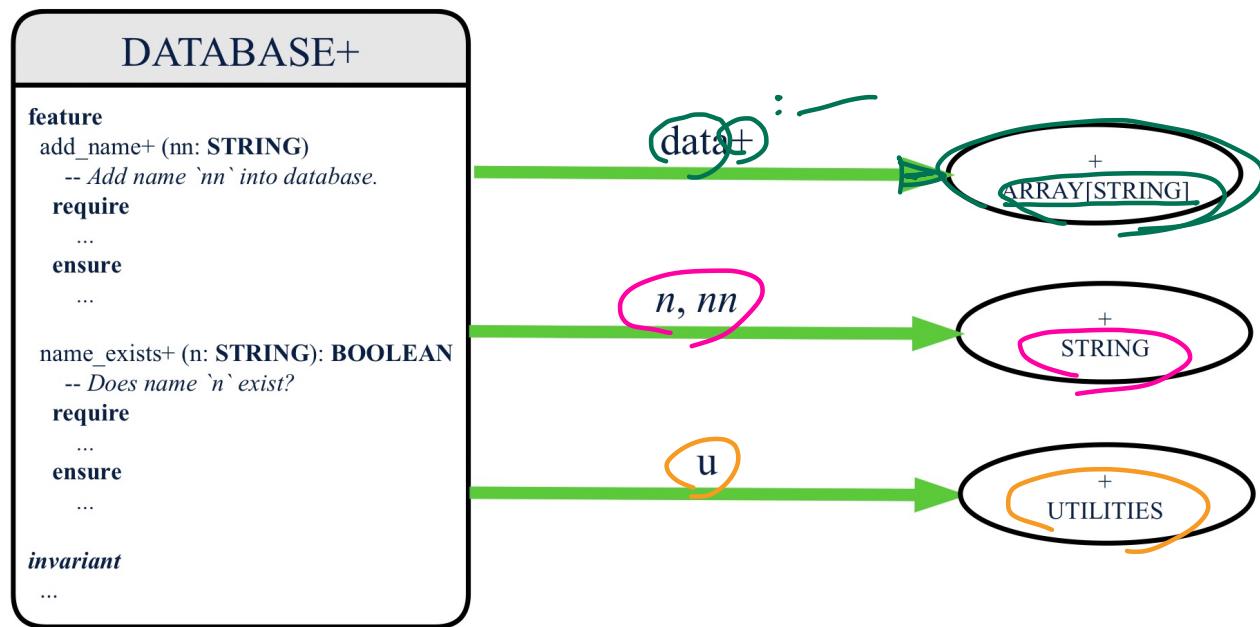
If STRING is to be emphasized, label is `data: ARRAY[...]`, where ... denotes the supplier class STRING being pointed to.



Presenting CS Relation in Diagram: Approach 2

If ARRAY is to be emphasized, label is `[data]`.

The supplier's name should be complete: ARRAY [STRING]



Programming Client-Supplier Relation

DESIGN ONE:

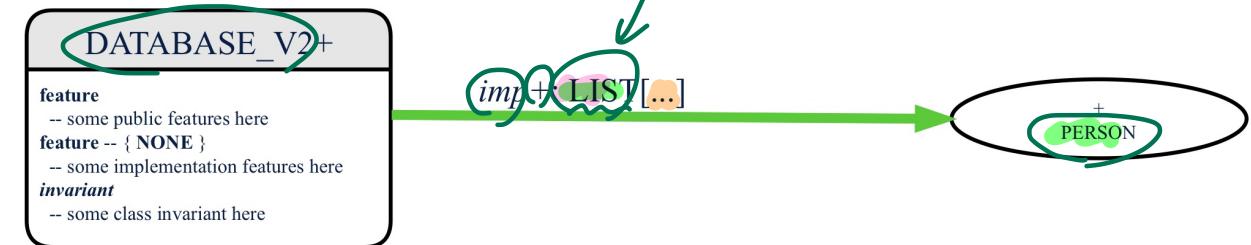
```
class DATABASE_V1
feature {NONE} -- implementation
    imp: ARRAYED_LIST[PERSON]
... -- more features and contracts
end
```

DESIGN TWO:

```
class DATABASE_V2
feature {NONE} -- implementation
    imp: LIST[PERSON]
... -- more features and contracts
end
```

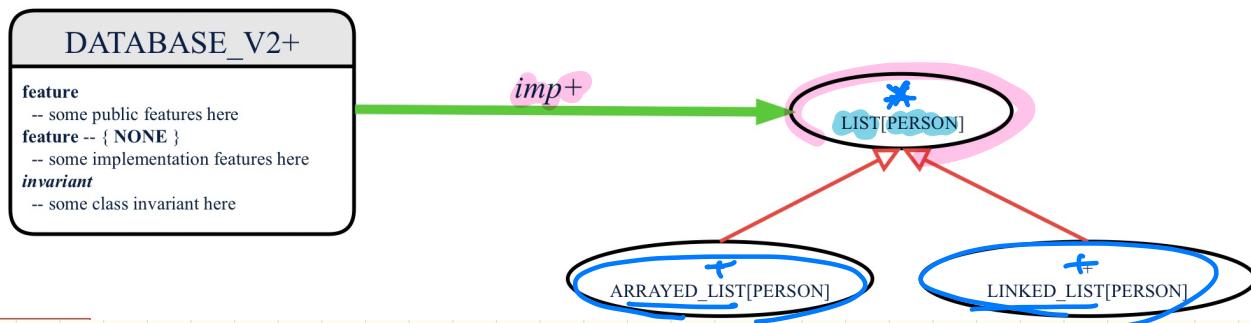
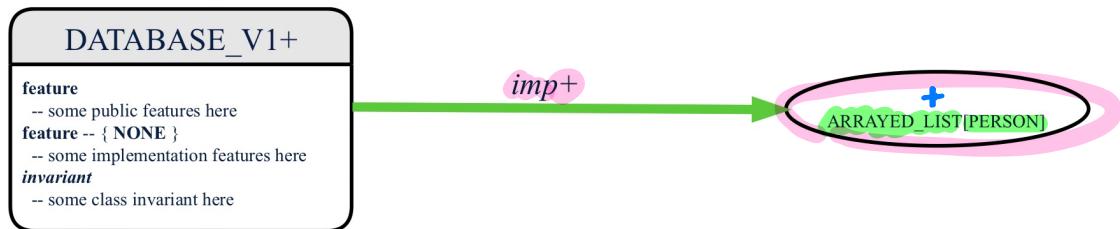
Presenting CS Relation in Diagram: Approach 1

We may focus on the PERSON supplier class, which may not help judge which design is better.



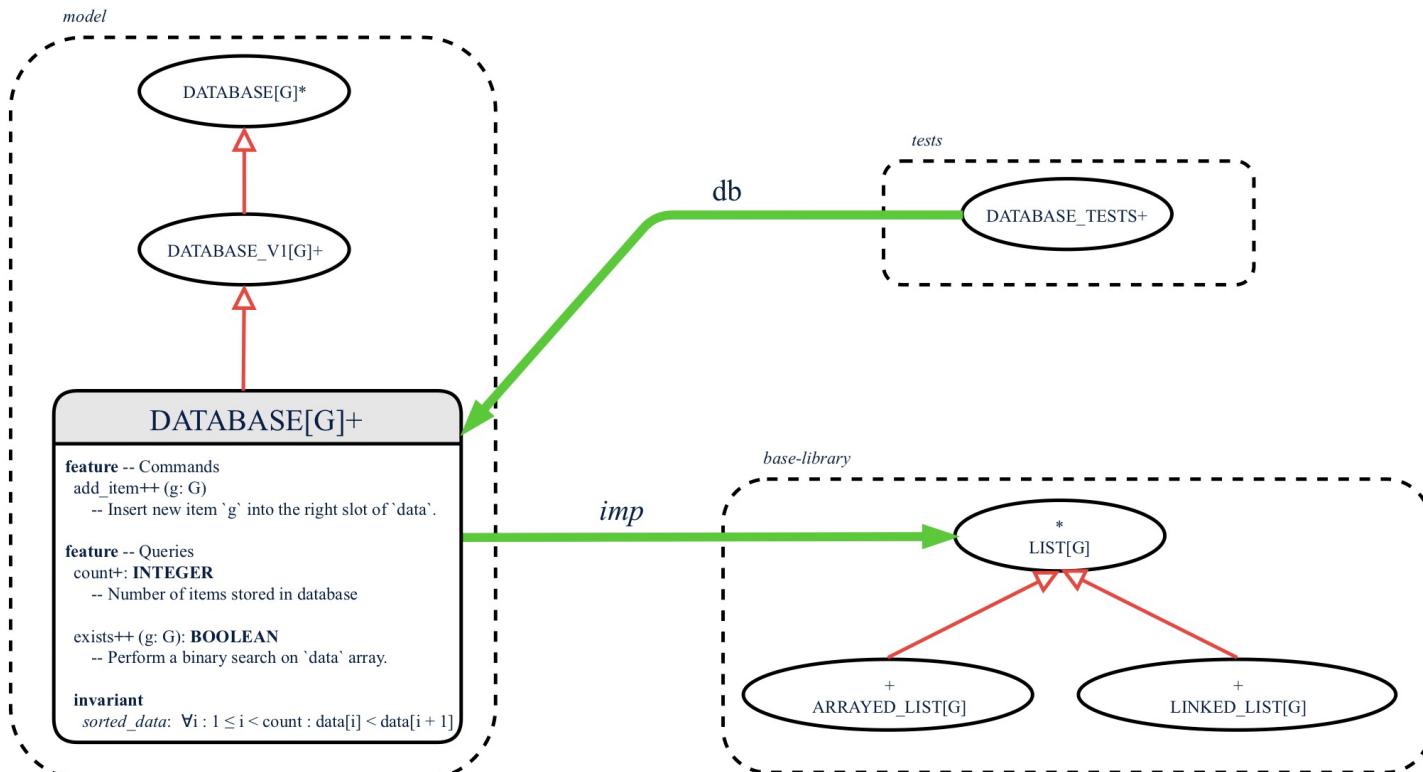
Presenting CS Relation in Diagram: Approach 2

Alternatively, we may focus on the LIST supplier class, which in this case helps us judge which design is better.



Clusters

Use *clusters* to group classes into logical units.



Lecture 4

Part 2

Postcondition: Asserting Set Equality

Writing Postcondition: Exercise

Problem

`all_positive_values (a: ARRAY[INTEGER]): ARRAY[INTEGER]`

require

`no_duplicates: ??`

ensure

`across Result is`

`all`

$x > 0$

`end`

incomplete

Witness/Counter-example

`all-p-vs (<< 2, 3, -1, 4, -2 >>)`



`<< 2, 3, 4 >>`

Correct

but the output
is incorrect

`all-p-vs (<< 2, 3, -1, 4, -2 >>)`

wrong
imp.

↳ `<< 1, 5, 7 >>`

↳ since each num. is $> 0 \Rightarrow$

↑ ↓
design is flawed!

Violation
no postcond.

Writing Postcondition: Exercise

`all_positive_values (a: ARRAY[INTEGER]): ARRAY[INTEGER]`

require

no_duplicates: ??

ensure

across Result is x

all

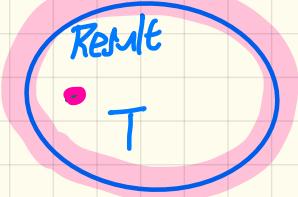
$x > 0$

end



Solution

`ARRAY[INTEGER]`



$$S \subseteq T \equiv (\forall x | [x \in S \Rightarrow x \in T] \quad S \subseteq T)$$

$$\quad \quad \quad (x \in T \Rightarrow x \in S) \quad T \subseteq S$$

`all_pos_in_a_also_in_result :`

across all x x is \neq a implies $Result.\text{has}(x)$ end-

`all_num_in_result_pos_and_`

`in_a :`

across x x is \neq a all $x > 0$ and $a.\text{has}(x)$ and

Lecture 4

Part 3

Abstraction of a Birthday Book

Testing REL in MATHMODELS

r.count → 9

Say $r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

- ✓ r.domain → {a, b, c, d, e, f} | r.image(a).count = 1
- r.range → {1, 2, 3, 4, 5, 6} ↳ singleton set
- ✓ r.image(a) → {1, 4} → functional domain → at most 1 value in range
- ✓ r[g] ↳ \emptyset ↳ domain value

Say $r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

$$\begin{aligned}
 & r \text{ overridden } \{(a, 3), (c, 4)\} \\
 & = \underbrace{\{(a, 3), (c, 4)\}}_t \cup \underbrace{\{(b, 2), (b, 5), (d, 1), (e, 2), (f, 3)\}}_{r \text{ domain subtracted } t \text{ domain}} \\
 & \qquad \qquad \qquad \{a, c\} \\
 & = \{(a, 3), (c, 4), (b, 2), (b, 5), (d, 1), (e, 2), (f, 3)\}
 \end{aligned}$$

(a, 3)
(c, 4)

Testing REL in MATHMODELS

Command : IMP. of model

Query : Contracts

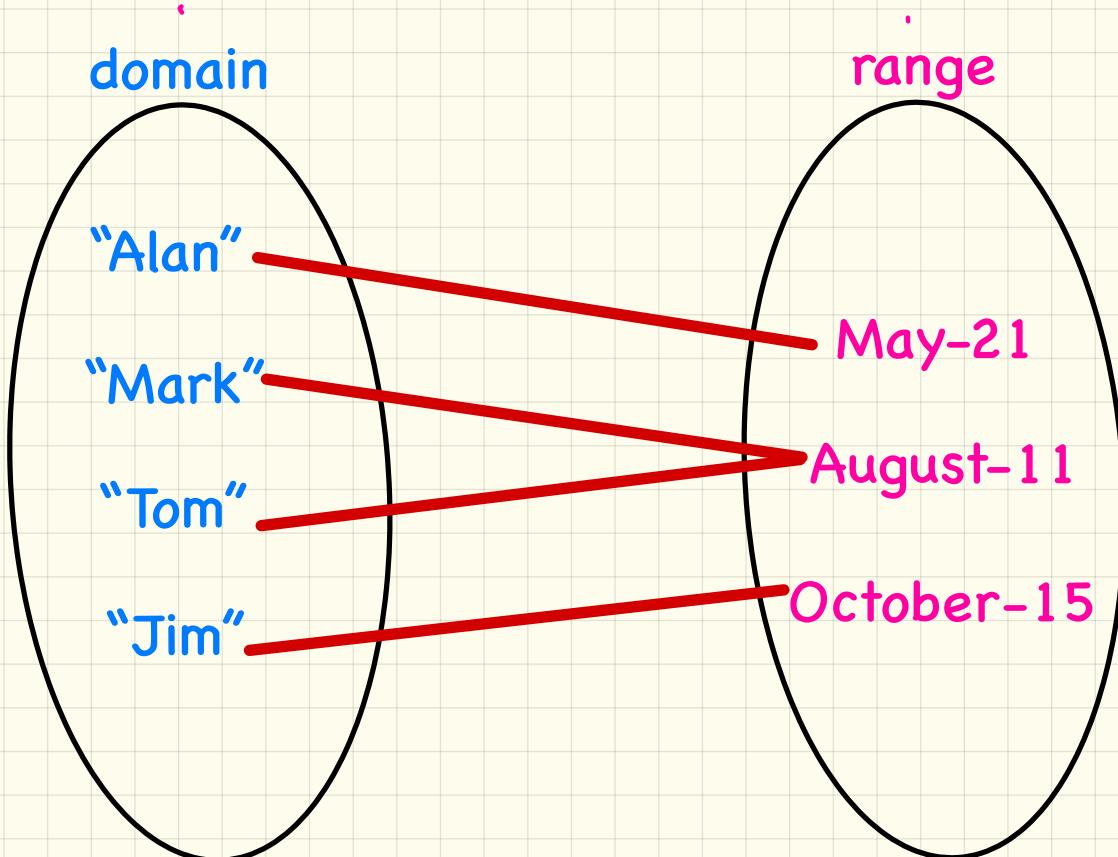
```
test_rel: BOOLEAN
local
  r, t: REL[STRING, INTEGER]
  ds: SET[STRING]
do
  create r.make_from_tuple_array (
    <<["a", 1], ["b", 2], ["c", 3],
    ["d", 4], ["b", 5], ["c", 6],
    ["d", 1], ["e", 2], ["f", 3]>>)
  create ds.make_from_array (<<"a">>)
-- r is not changed by the query 'domain_subtracted'
  t := r.domain_subtracted (ds)
Result :=
  t ~ r and not t.domain.has ("a") and r.domain.has ("a")
check Result end
-- r is changed by the command 'domain_subtract'
  r.domain_subtract (ds)
Result :=
  t ~ r and not t.domain.has ("a") and not r.domain.has ("a")
end
```



Say $r = \{(a, 1), (b, 2), (c, 3), (a, 4), (b, 5), (c, 6), (d, 1), (e, 2), (f, 3)\}$

- **r.domain**: set of first-elements from r
 - $r.domain = \{ d \mid (d, r) \in r \}$
 - e.g., $r.domain = \{a, b, c, d, e, f\}$
- **r.range**: set of second-elements from r
 - $r.range = \{ r \mid (d, r) \in r \}$
 - e.g., $r.range = \{1, 2, 3, 4, 5, 6\}$
- **r.inverse**: a relation like r except elements are in reverse order
 - $r.inverse = \{ (r, d) \mid (d, r) \in r \}$
 - e.g., $r.inverse = \{(1, a), (2, b), (3, c), (4, a), (5, b), (6, c), (1, d), (2, e), (3, f)\}$
- **r.domain_restricted(ds)**: sub-relation of r with domain ds .
 - $r.domain_restricted(ds) = \{ (d, r) \mid (d, r) \in r \wedge d \in ds \}$
 - e.g., $r.domain_restricted(\{a, b\}) = \{\text{(a, 1), (b, 2), (a, 4), (b, 5)}\}$
- **r.domain_subtracted(ds)**: sub-relation of r with domain not ds .
 - $r.domain_subtracted(ds) = \{ (d, r) \mid (d, r) \in r \wedge d \notin ds \}$
 - e.g., $r.domain_subtracted(\{a, b\}) = \{\text{(c, 6), (d, 1), (e, 2), (f, 3)}\}$
- **r.range_restricted(rs)**: sub-relation of r with range rs .
 - $r.range_restricted(rs) = \{ (d, r) \mid (d, r) \in r \wedge r \in rs \}$
 - e.g., $r.range_restricted(\{1, 2\}) = \{\text{(a, 1), (b, 2), (d, 1), (e, 2)}\}$
- **r.range_subtracted(ds)**: sub-relation of r with range not ds .
 - $r.range_subtracted(rs) = \{ (d, r) \mid (d, r) \in r \wedge r \notin rs \}$
 - e.g., $r.range_subtracted(\{1, 2\}) = \{\text{(c, 3), (a, 4), (b, 5), (c, 6)}\}$

Model of an Example Birthday Book



Birthday Book: Design

BIRTHDAY_BOOK

model: FUN[NAME, BIRTHDAY]
-- abstraction function

count: INTEGER
-- number of entries

put(NAME, d:BIRTHDAY)
ensure

model_operation: model ~ (old model.deep_twin) overridden_by (n,d)
-- infix symbol for override operator: @<+

remind(d: BIRTHDAY) ARRAY[NAME]
ensure

nothing_changed: model ~ (old model.deep_twin)

same_counts: Result.count = (model.range_restricted_by(d)).count

same_contents: $\forall \text{ name} \in (\text{model.range_restricted_by}(d)).\text{domain}$: name $\in \text{Result}$

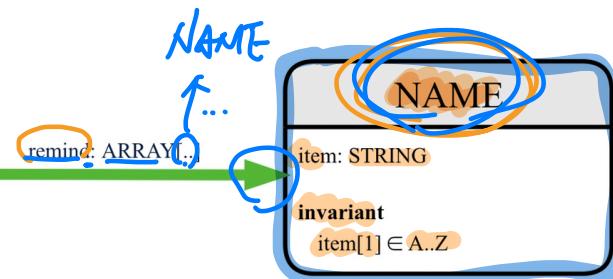
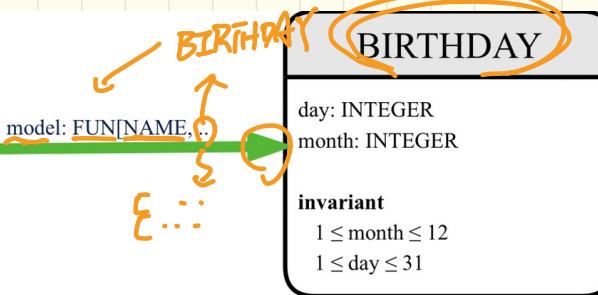
-- infix symbol for range restriction: model @> (d)

invariant:

consistent_book_and_model_counts: count = model.count

model ~ (old model.d-t) @<+ [n,d]

guard

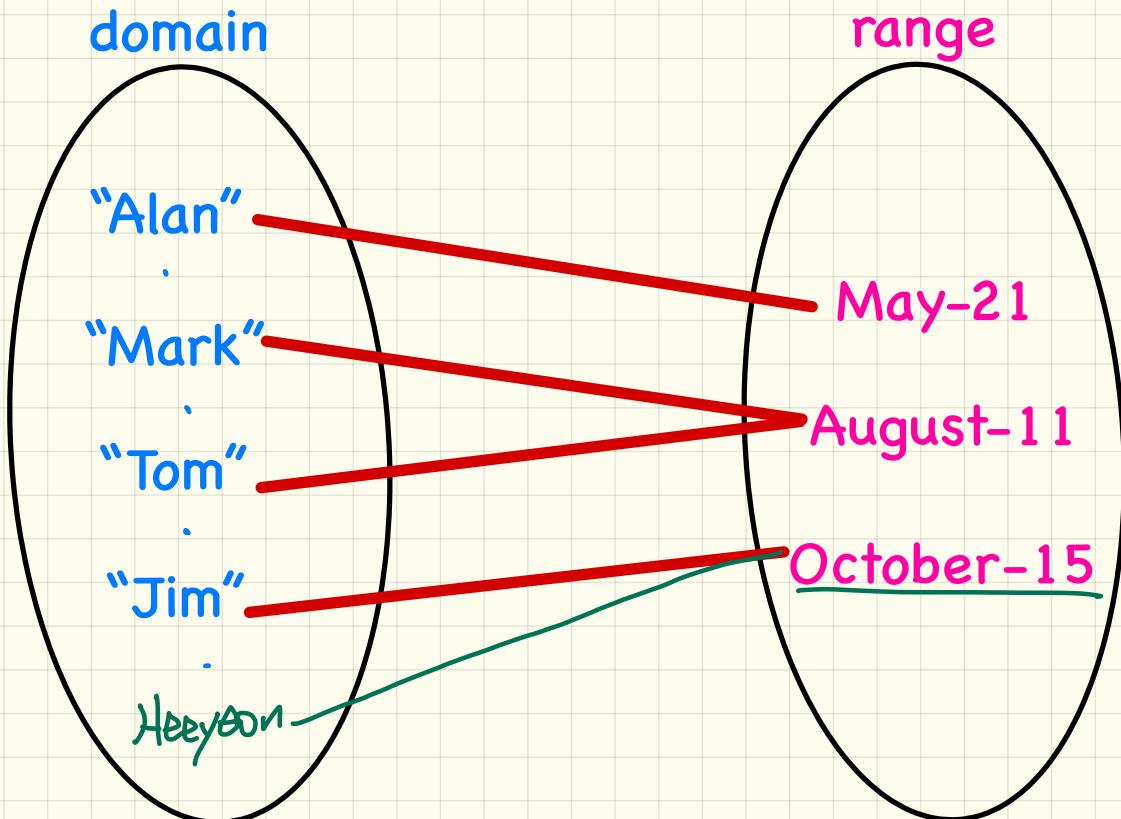


STRING

"12 Alan"
↓
STRING

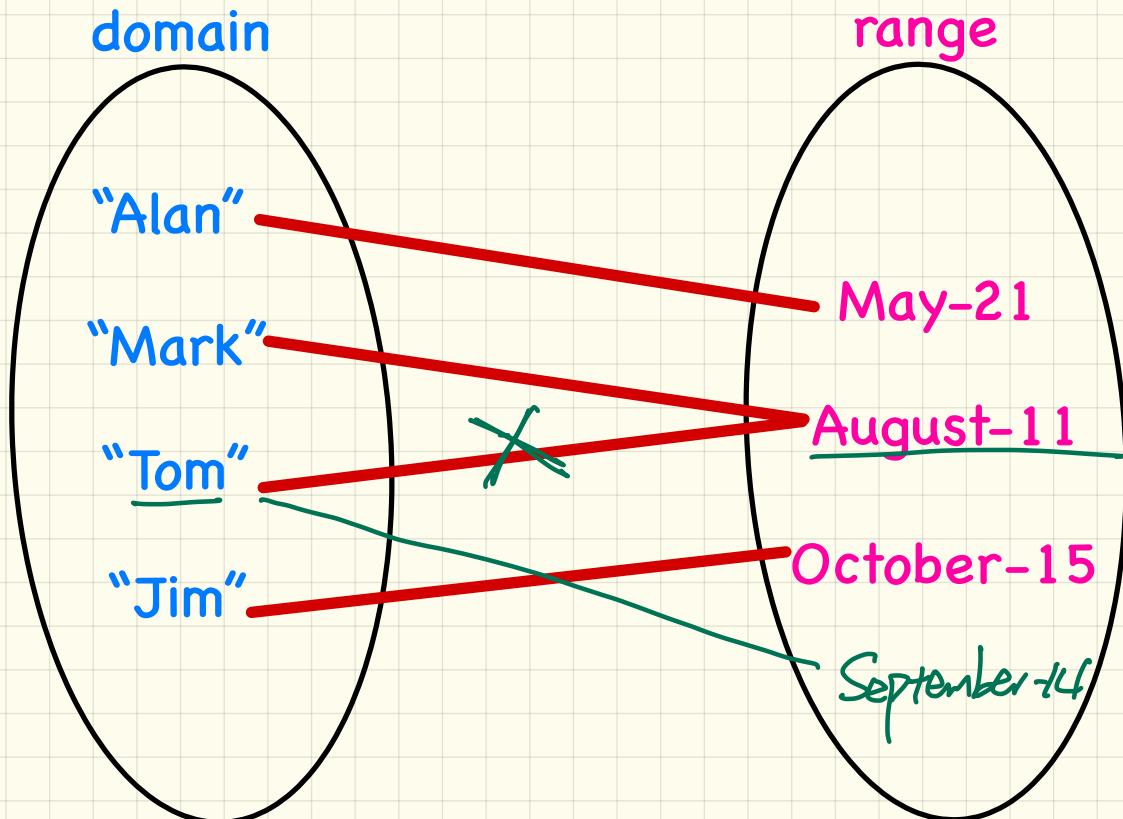
Birthday Book: Model Operation (1.1)

book.put("Heeyeon", October-15)



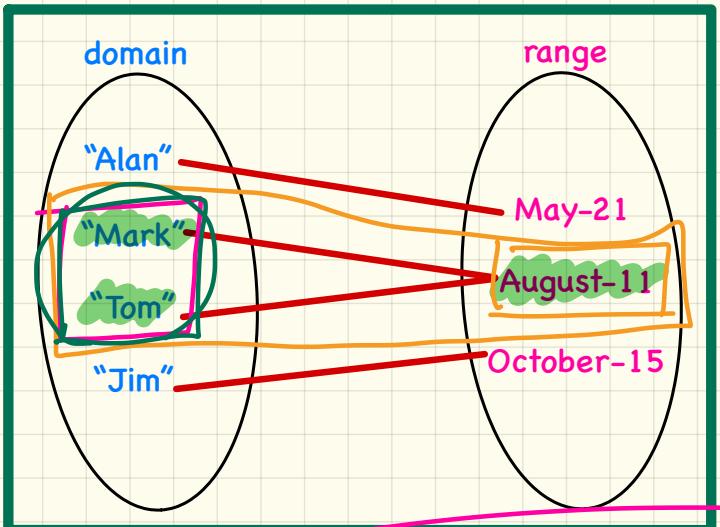
Birthday Book: Model Operation (1.2)

book.put("Tom", September-14)



Birthday Book: Model Operation (2.1)

book.remind(August-11)

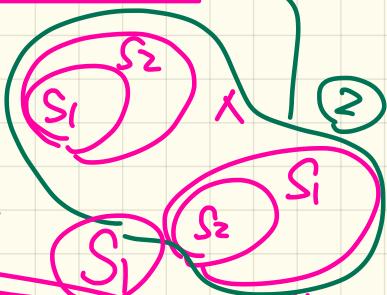


$$S_1 = S_2$$

$S_1 \subseteq S_2$
 $S_2 \subseteq S_1$

$$\hookrightarrow$$

$S_1 \subseteq S_2$
 $S_2 \subseteq S_1$



(model range-restricted by (August-11)). domain

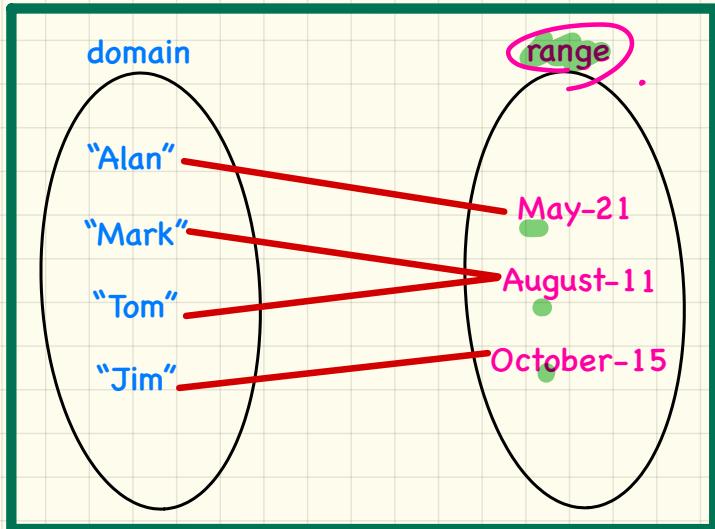
relation

~ Result Array
S2

Birthday Book: Model Operation (2.2)

book.remind(November-29)

→ ~~empty array~~.



Birthday Book: Implementation

BIRTHDAY_BOOK

```
model: FUN[NAME, BIRTHDAY]
-- abstraction function
do
  -- promote hashtable to function
  ensure
    same_counts: Result.count = implementation.count
    same_contents: ∀ [name, date] ∈ Result: [name, date] ∈ implementation
end

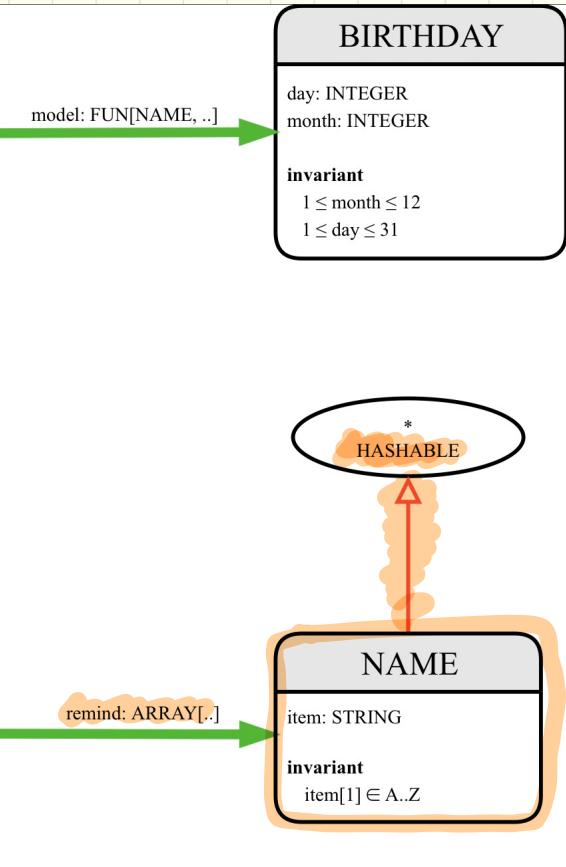
put(n: NAME; d: BIRTHDAY)
do
  -- implement using hashtable
  ensure
    model_operation: model ~ (old model.deep_twin) @<+ [n,d]
end

remind(d: BIRTHDAY): ARRAY[NAME]
do
  -- implement using hashtable
  ensure
    nothing_changed: model ~ (old model.deep_twin)
    same_counts: Result.count = (model @> d).count
    same_contents: ∀ name ∈ (model @> d).domain: name ∈ Result
end

count: INTEGER -- number of names

feature {NONE}
implementation: HASH_TABLE[BIRTHDAY, NAME]

invariant:
  consistent_book_and_model_counts: count = model.count
  consistent_book_and_imp_counts: count = implementation.count
```



Lecture 4

Part 4

Design Pattern: Iterator

Principle of Information Hiding

Supplier:

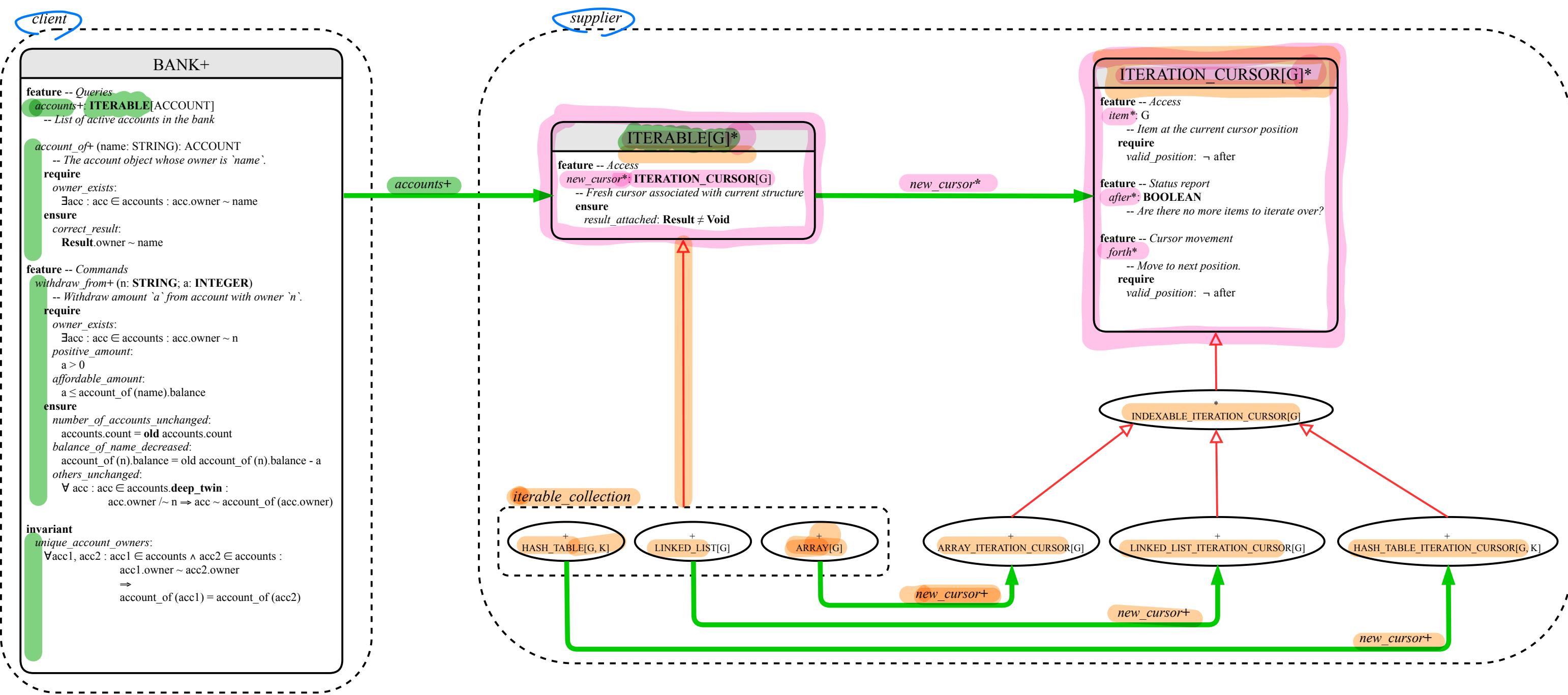
```
class
  CART
feature
  orders: ARRAY[ORDER]
end

class
  ORDER
feature
  price: INTEGER
  quantity: INTEGER
end
```

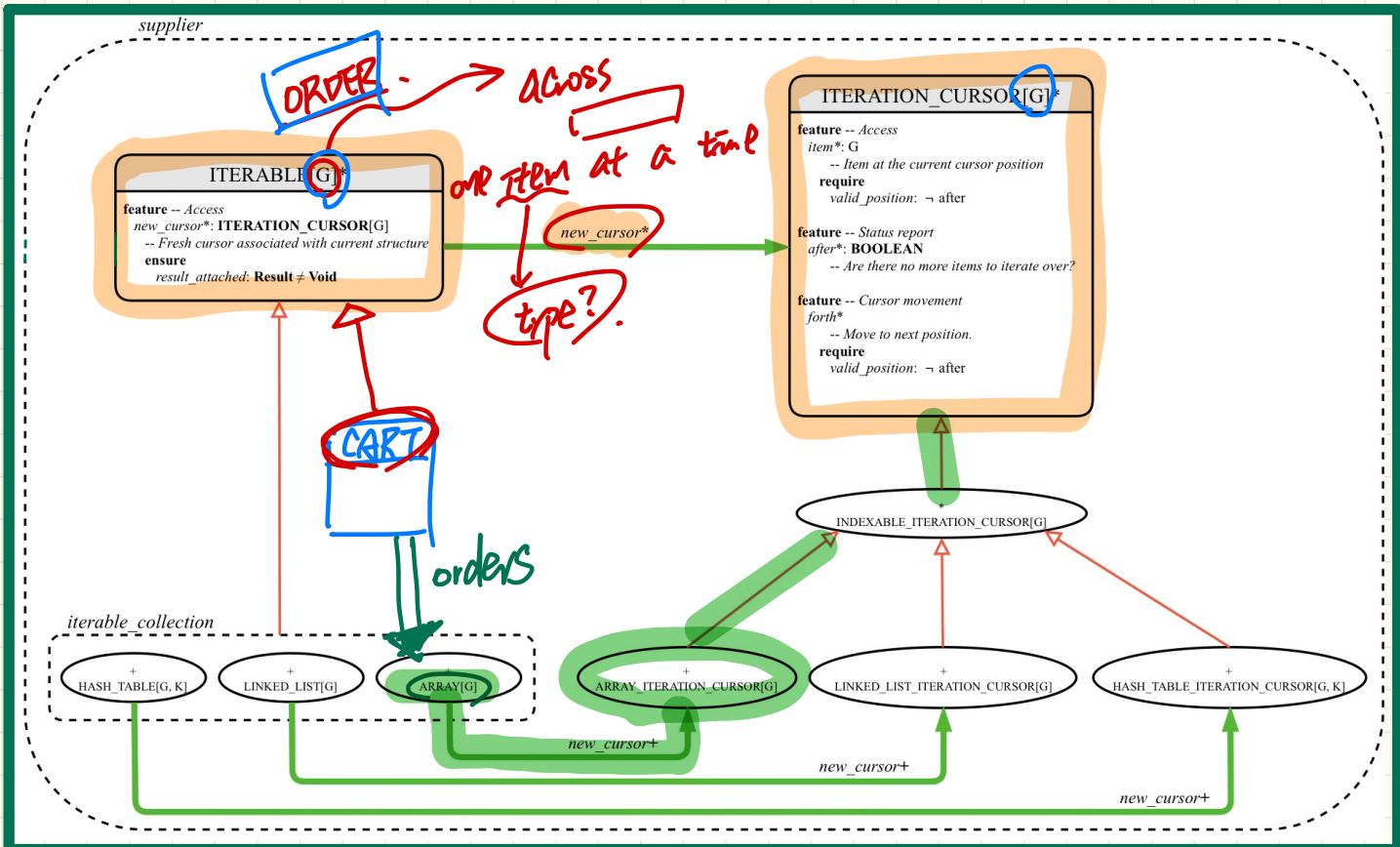
Problems?

Client:

```
class
  SHOP
feature
  cart: CART
  checkout: INTEGER
  do
    from
      i := cart.orders.lower
    until
      i > cart.orders.upper
    do
      Result := Result +
        cart.orders[i].price
      *
      cart.orders[i].quantity
      i := i + 1
    end
  end
end
```



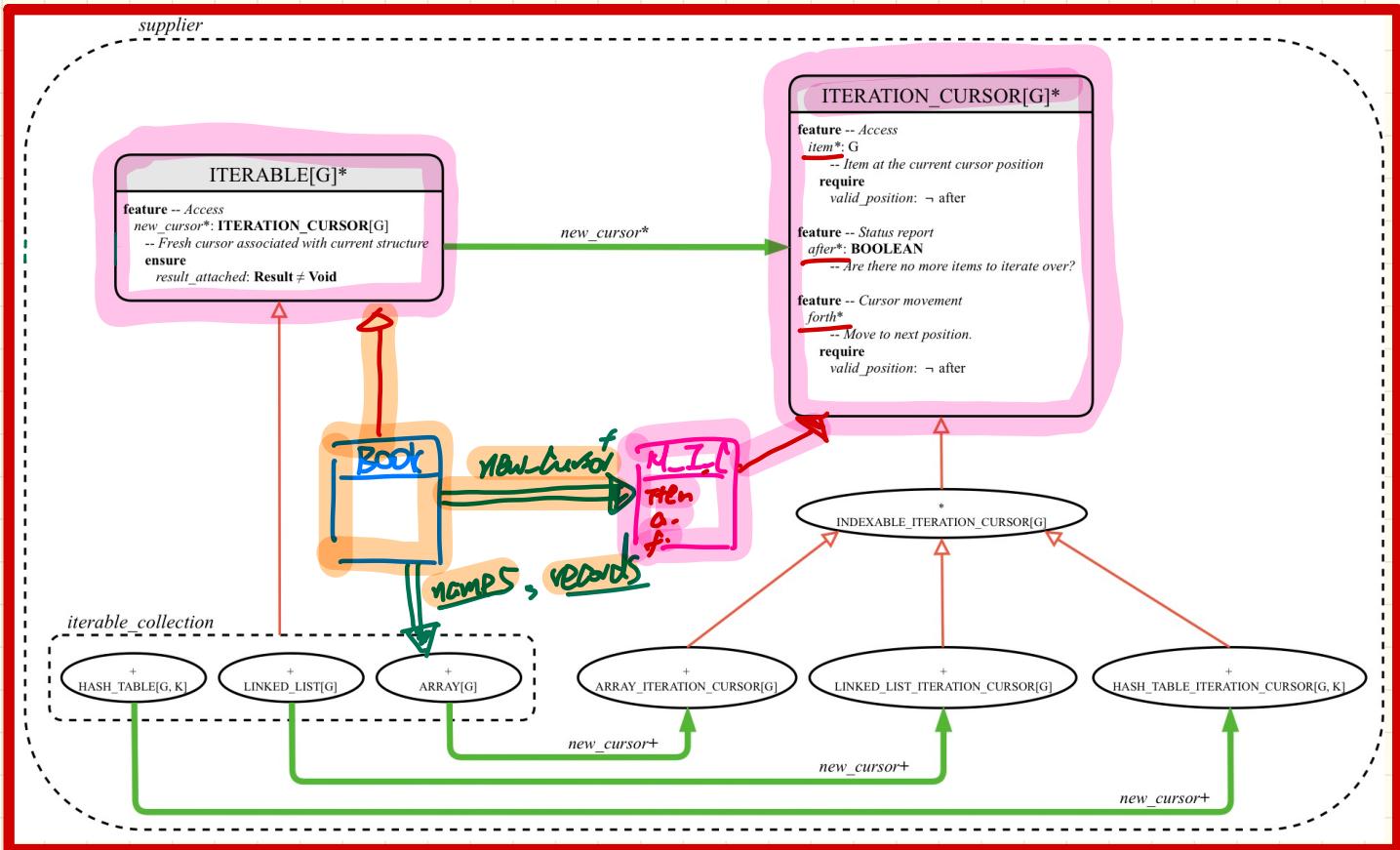
Implementing the Iterator Pattern: Easy Case



Implementing the Iterator Pattern: Easy Case

```
class CART
  inherit ITERABLE[ORDER]
  feature {NONE}
    orders: ARRAY[ORDER]
  featureP --
    new_cursor : I-List[ORDER]
    do
      Result := orders.new_cursor
    end
  end
```

Implementing the Iterator Pattern: Hard Case



Implementing the Iterator Pattern: Hard Case

class

BOOK[G]

data

Inherit

ITERABLE[TUPLE[STRING, G]]

feature {NONE}

names: ARRAY[STRING]

records: ARRAY[G]

feature --

new_cursor :

I_C[TUPLE[S,G]]

do

create {M-I-C[S,G]}.

make(-->).

end

across

bb

private

names

(
"ahit")

--

(
"Dct-fb")

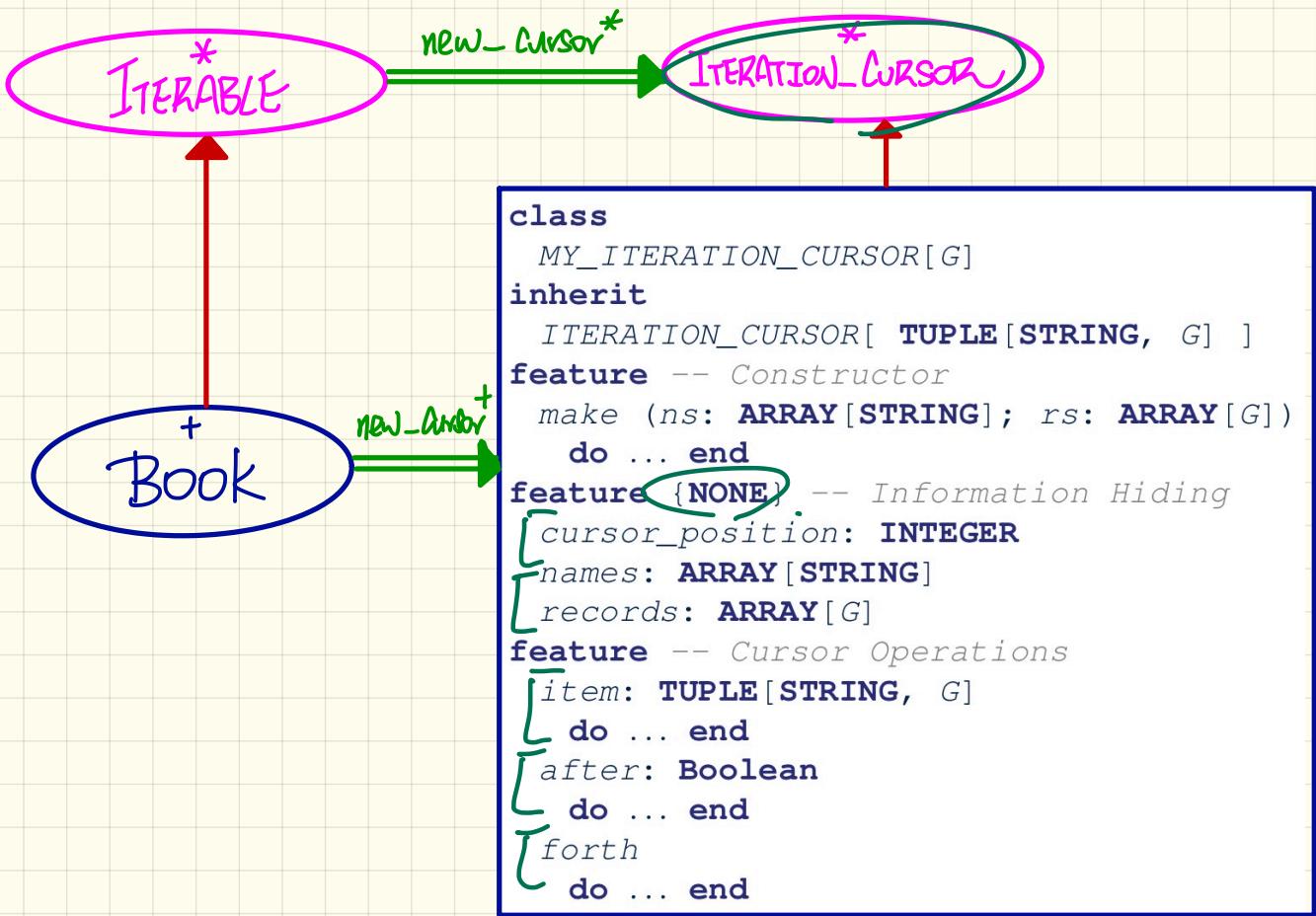
--

(
"records")

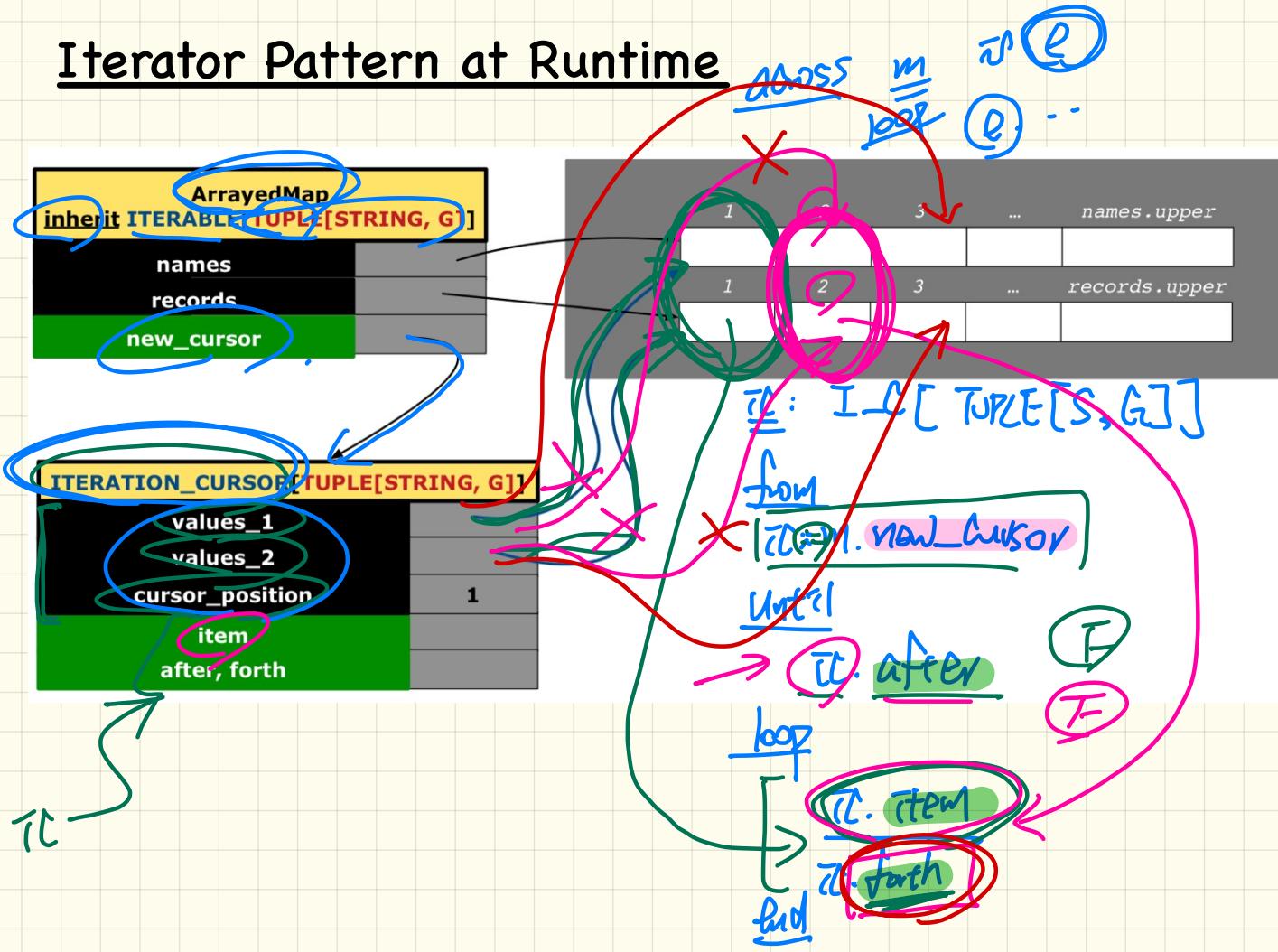
["ahit", off-10]

TUPLE[--, --]

Implementing the Iterator Pattern: Hard Case



Iterator Pattern at Runtime



Use of Iterable in Contracts

```
class CHECKER .  
feature -- Attributes  
  collection: ITERABLE [INTEGER]  
feature -- Queries  
  is_all_positive: BOOLEAN  
    -- Are all items in collection positive?  
  do  
    ...  
  ensure  
    across  
      collection is item  
      all  
        item > 0  
      end  
    end
```

does not support [...]

```
class BANK  
  ...  
  accounts: LIST [ACCOUNT]  
  binary_search (acc_id: INTEGER): ACCOUNT  
    -- Search on accounts sorted in non-descending order.  
  require  
    across  
      1 | ... | (accounts.count - 1) is i  
      all  
        accounts[i].id <= accounts[i + 1].id  
      end  
    do  
      ...  
    ensure  
      Result.id = acc_id  
    end
```

ITERABLE X

Item Item
↳ specific to LIST

Use of Iterable in Contracts: Exercise

```
class BANK
...
accounts: LIST [ACCOUNT]
contains_duplicate: BOOLEAN
    -- Does the account list contain duplicate?
do
...
ensure
     $\exists i, j : \text{INTEGER} \mid$ 
     $1 \leq i \leq \text{accounts.count} \wedge 1 \leq j \leq \text{accounts.count}$ 
     $\text{accounts}[i] \sim \text{accounts}[j] \Rightarrow i = j$ 
end
```

across

Use of Iterable in Implementation (1)

```
class BANK
  accounts: ITERABLE[ACCOUNT]
  max_balance: ACCOUNT
    -- Account with the maximum balance value.
  require ???
  local
    cursor: ITERATION_CURSOR[ACCOUNT]; max: ACCOUNT
  do
    from cursor := accounts.new_cursor; max := cursor.item
    until cursor.after
  loop do
    if cursor.item.balance > max.balance then
      max := cursor.item
    end
    cursor.forth
  end
  ensure ???
end
```

use from-until loop
if you wish to work with
an I-C.

Use of Iterable in Implementation (2)

```
class SHOP
  cart: CART
  checkout: INTEGER
    -- Total price calculated based on orders in the cart.
  require ???
  do
    across
      cart is order
    loop
    Result := Result + order.price * order.quantity
  end
  ensure ???
end
```

Annotations:

- Yellow box around `???` in `require`.
- Yellow box around `cart` in `cart is order`.
- Green box around `order` in `order.price` and `order.quantity`.
- Handwritten note: $\Rightarrow \text{rc} := \text{Cart.new-cursor}$
- Handwritten note: rc.item.
- Handwritten note: *rc forth*

ITERABLE ?

```
class BANK
  accounts: LIST[ACCOUNT] -- Q: Can ITERABLE[ACCOUNT] work?
  max_balance: ACCOUNT
    -- Account with the maximum balance value.
  require ???
  local
    max: ACCOUNT
  do
    max := accounts [1]
    across
      accounts is acc
    loop
      if acc.balance > max.balance then
        max := acc
      end
    end
  ensure ???
end
```

Annotations:

- Red X over `LIST[ACCOUNT]` in `accounts`.
- Yellow box around `???` in `require`.
- Yellow box around `accounts` in `accounts is acc`.
- Yellow box around `acc` in `acc.balance`.
- Yellow box around `max` in `max := acc`.
- Yellow box around `???` in `ensure`.

Lecture 4

Part 5

Exercise: Generics in Iterator

Exercise

```
defered class
  ITERABLE [G]
feature -- Access
  new_cursor: ITERATION_CURSOR [G]
  deferred end
end
```

new_cursor*

```
defered class
  ITERATION_CURSOR [G]
feature -- Cursor features
  item: G
  deferred end
  after: BOOLEAN
  deferred end
  forth
  deferred end
```

```
test_database: BOOLEAN
local
  db: DATABASE[STRING, INTEGER]
  tuples: LINKED_LIST[TUPLE[INTEGER, STRING]]
do
  create db.make
  create tuples.make
  across
    db is t
  loop
    tuples.extend (t)
  end
end
```

qp +

```
class
  DATABASE[G, H]
inherit
  ITERABLE [ ]
feature {NONE} -- Implementation
  gs: ARRAY[G]
  hs: ARRAY[H]
feature -- Iterable
  new_cursor: ITERATION_CURSOR[ ]
  local
    db_cursor: ITEM_ITERATION_CURSOR[H, G]
  do
    create db_cursor.make ( )
    Result := db_cursor
  end
end
```

new_cursor+

```
class
  ITEM_ITERATION_CURSOR[M, N]
inherit
  ITERATION_CURSOR[ ]
create
  make
feature {NONE} -- Implementation
  ms: ARRAY[M]
  ns: ARRAY[N]
feature -- Constructor
  make (new_ns: ARRAY[N]; new_ms: ARRAY[M])
    do ... end
feature -- Cursor features
  item: 
    do ... end
  after: BOOLEAN
    do ... end
  forth
    do ... end
end
```